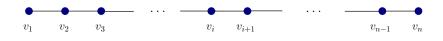
# **Dynamic Programming**

- (Weighted) Independent Set in Graphs
- Weighted Independent Sets in Path
- Dynamic Programming Formulation
- Implementation and Backtracking

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## The Path Graph

The path graph is a connected graph with two nodes of degree 1 and the other n-2 vertices of degree 2



Number of edges = 
$$\frac{1 + 2(n-2) + 1}{2} = n-1$$

So a path is a tree

**Input:** A node weighted graph  $G = (V, E), w \ w : V \to \mathbb{R}^+$ 

Output: An independent set of G of maximum cardinality weight

A company wants to open restaurants on the motorway

- Designated service areas  $s_1, ..., s_n$  every 7 kilometers
- A restaurant at  $s_i$  gives estimated profit  $p_i$
- No two restaurants can be located within 10 km of each other

#### Select a subset of sites to maximize total profit

Problem can be modeled by a node weighted path graph

- Each site  $s_i$  is a vertex with weight equal to  $p_i$
- If two sites are within 10 km of each other make an edge between the corresponding vertices 

  ▷ note: we get a path graph

**Input:** A node weighted path graph P = (V, E),  $w : V \to \mathbb{R}^+$ 

**Output:** An independent set of *P* of maximum weight

No consecutive vertices can be chosen



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An independent set of weight 16

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An independent set of weight 22

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### **Greedy Approach:**

- Select a node with max weight
- Mark its neighbors as incompatible
- Repeat the process with remaining unmarked nodes

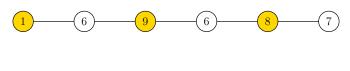


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#### Divide & Conquer approach-1:

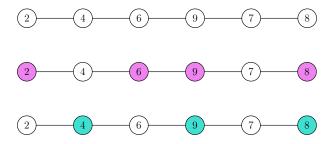
- Divide *P* into left and right halves
- Find max weight independent sets in both
- Combine the two sets to get the answer



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#### Divide & Conquer approach-2:

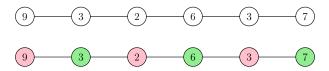
- Divide P into odd and even indexed vertices
- Each one is an independent set
- Return the larger of the two



**Input:** A node weighted path graph P = (V, E),  $w : V \to \mathbb{R}^+$ **Output:** An independent set of P of maximum weight

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