Polynomial Time Reduction

- Polynomial Time Reduction Definition
- Reduction by Equivalence
- Reduction from Special Cases to General Case
- Reduction by Encoding with Gadgets
- Transitivity of Reductions
- Decision, Search and Optimization Problem
- Self-Reducibility

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Transitivity of Reductions

Problem A is polynomial time reducible to Problem B, $A \leq_p B$

If any instance of problem A can be solved using a polynomial amount of computation plus a polynomial number of calls to a solution of problem B

We used the following techniques for reduction

- Simple Equivalence
- Special Case to General Case
- Encoding with Gadgets

A very powerful technique is to exploit transitivity of reductions

Theorem: If $X \leq_p Y$ and $Y \leq_p Z$, then $X \leq_p Z$

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Theorem: If $X \leq_p Y$ and $Y \leq_p Z$, then $X \leq_p Z$

Proof: Let $\mathcal{A}_Z$ be an algorithm for Z

Given any instance I_X of X we will solve X on I_X using $\mathcal{A}_Z^+$

- There is an algorithm $\mathcal{A}_Y$ for Y using $\mathcal{A}_Z^+$ (maybe many others too)
- There is an algorithm $\mathcal{A}_X$ for X using $\mathcal{A}_Y^+$

$\mathcal{B}_\mathcal{X}$: the new algorithm for X uses everything as of $\mathcal{A}_X$ but it uses the specific algorithm $\mathcal{A}_Y$ that is built upon $\mathcal{A}_Z$

▷ $\mathcal{B}_\mathcal{X}$ essentially composes the two reductions into one

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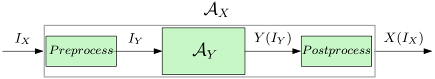

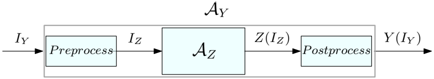

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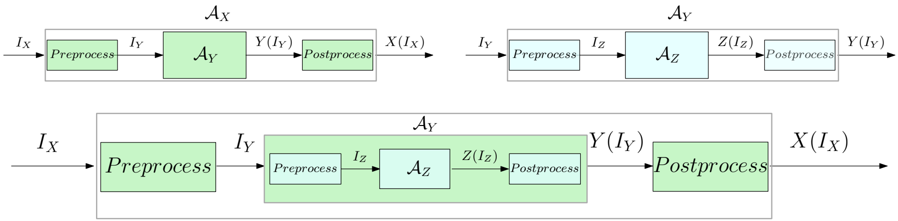

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Transitivity is an extremely useful property of reduction

- $\textsc{sat}(f) \leq_p 3\text{-}\textsc{sat}(f')$ and $3\text{-}\textsc{sat}(f) \leq_p \textsc{ind-set}(G, k)$
  - From these we conclude that $\textsc{sat}(f) \leq_p \textsc{ind-set}(G, k)$
- $\textsc{sat}(f) \leq_p 3\text{-}\textsc{sat}(f') \leq_p \textsc{ind-set}(G, k) \leq_p \textsc{vertex-cover}(G, t) \leq_p \textsc{set-cover}(U, \mathcal{S}, l)$
  - From these we conclude that $\textsc{sat}(f) \leq_p \textsc{set-cover}(U, \mathcal{S}, l)$
- And many others