## Intractable Problems

- Clique
- Independent Set
- Vertex Cover
- Set Cover
- Set Packing
- Satisfiability Problem
- Hamiltonian Cycle and Path

- Traveling Salesman Problem
- Graph Coloring
- Circuit Satisfiability
- Knapsack
- Subset Sum
- Prime and Factor
- Partition

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# Vertex Cover

An vertex cover in a graph is subset C of vertices such that each edge has at least one endpoint in C



The VERTEX-COVER(G, k) problem: Is there a cover of size k in G?

## Network Security: Rout Based Filtering

- Identify a small set of routers/AS
- So as all packets can be monitored at those routers/switches
  - check if the source/destination addresses are valid given the routing table and network topology
- Route-based distributed packet filtering
- Prevents distributed denial of service attacks

## Set Cover

• Given a set U of n elements

• A collection S of m subsets  $S_1, S_2, \ldots, S_m \subseteq U$ 

• A Set Cover is a subcollection  $I \subset \{1, 2, ..., m\}$  with  $\bigcup_{i \in I} S_i = U$ 

$$\begin{split} & \mathcal{U} := \{1,2,3,4,5,6\} \\ & \mathcal{S} := \big\{\{1,2,3\},\{3,4\},\{1,3,4,5\},\{2,4,6\},\{1,3,5,6\},\{1,2,4,5,6\}\big\} \\ & \text{Cover-1: } \{1,2,3\}, & \{1,3,4,5\},\{2,4,6\} \\ & \text{Cover-2: } \{1,2,3\}, & \{1,3,4,5\}, & \{1,2,4,5,6\} \\ & \text{Cover-3: } & \{1,3,4,5\}, & \{1,2,4,5,6\} \end{split}$$

Cover-1 has size 3, the latter two have size 2 each

## Set Cover

- Given a set U of n elements
- A collection  $\mathcal S$  of m subsets  $S_1, S_2, \ldots, S_m \subseteq U$
- A Set Cover is a subcollection  $I \subset \{1, 2, ..., m\}$  with  $\bigcup_{i \in I} S_i = U$



The SET-COVER(U, S, k) problem: Is there a cover of size k for U?

#### System Integration

- *U* is the set of capabilities we want our system to have
- S: Available softwares each providing a **subset** of capabilities
- Select a (small) subset of softwares to provide all functionalities

#### IBM antivirus tool

- U is the set of (500) known viruses
- They found a set of about 9000 strings of 20 bytes or more that occur in the binaries of viruses but not in "clean" codes
- $\blacksquare\ \mathcal{S}$  : For each string the subset of viruses containing it
- Select a (small) subset of strings that should be searched in codes to detect any virus

# Set Packing

- Given a set U of n elements
- A collection  $\mathcal{S}$  of m subsets  $S_1, S_2, \ldots, S_m \subseteq U$
- A subcollection  $I \subset \{1, 2, ..., m\}$  pack together if for all  $i \neq j \in I$  $S_i \cap S_j = \emptyset$

$$U := \{1, 2, 3, 4, 5, 6\}$$
  
$$S := \{\{1, 2, 3\}, \{4, 5\}, \{4, 6\}, \{2, 3\}, \{1, 6\}, \{4, 5, 6\}\}$$

Pack-1:  $\{1, 2, 3\}, \{4, 5\}$ Pack-2:  $\{4, 5\}, \{2, 3\}, \{1, 6\}$ Pack-3:  $\{1, 2, 3\}, \{4, 5, 6\}$ 

Pack-1 and Pack-3 have size 2 each, Pack-3 has size 3

# Set Packing

- Given a set U of n elements
- A collection S of m subsets  $S_1, S_2, \ldots, S_m \subseteq U$
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The SET-PACKING (U, S, k) problem: Is there a packing of size k?

#### **Resource Sharing**

- *U* is a set of non sharable resources
- $\mathcal{S}$  : Processes each requesting a subset of resources
- Select a (large) subset of processes to allocate resources to

## Airline Crew Scheduling

- U is the set of crew staff
- $\mathcal{S}$  : teams of members willing to work with each other
- Select (many) teams to service a set of flights