Randomized Computation

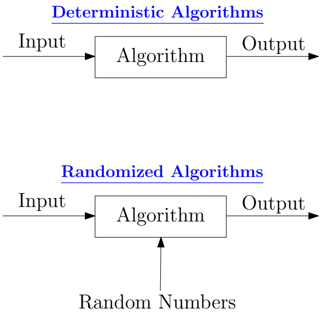
- Deterministic and (Las Vegas & Monte Carlo) Randomized Algorithms
- Probability Review
- Probabilistic Analysis of deterministic QUICK-SORT Algorithm
- RANDOMIZED-SELECT and RANDOMIZED-QUICK-SORT
- Max-Cut
- Min-Cut
- MAX-3-SAT and Derandomization
- Closest Pair
- Randomized Complexity Classes

IMDAD ULLAH KHAN

Algorithmic Design Paradigm

We have seen **Deterministic** algorithms using

- Greedy
- Divide-and-Conquer
- Dynamic Programming
- Network Flows



Randomized Algorithms incorporate randomness in their operation

Randomized Algorithm

Randomized Algorithm receives, in addition to the input, a random number stream to make random decision during execution

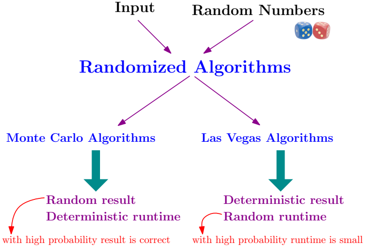
May give different results on the same input in different runs

Often aims for properties like:

- Good average-case (expected) behavior
- Getting exact answers with high probability
- Getting answer close to the right answer with high probability
- Runtime is small with high probability

Advantages:

- Simple and elegant. Their output/runtime is good with high probability
- The execution time or space requirement is smaller than that of the best deterministic algorithm



- **Monte Carlo Algorithms**
  - Guaranteed to run in a fixed time
  - Outputs a correct answer with some (high) probability
  - e.g. Min cut algorithm
- **Las Vegas Algorithms**
  - Guaranteed to output the correct answer
  - Running time is random (with high probability runtime is small)
  - e.g: randomized Quicksort, Closest pair

Las Vegas and Monte Carlo Algorithms

Input: An array A with $n/4$ 1's and $3n/4$ 0's

Output: An index k such that $A[k] = 1$

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
0	0	0	1	0	0	0	0	0	0	1	0	0	1	1	0

Algorithm Deterministic

```
k ← 0
for i = 1 → n do
    if A[i] = 1 then
        k ← i
        return k
```

Quality: correct

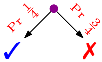
worst case runtime: $\frac{3n}{4}$

Algorithm Monte Carlo

```
k ← RANDOM(1 ··· n)
return k
```

Quality: correct w.p $\frac{1}{4}$

worst case runtime: 1

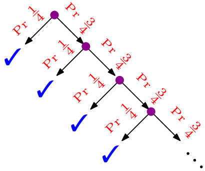


Algorithm Las Vegas

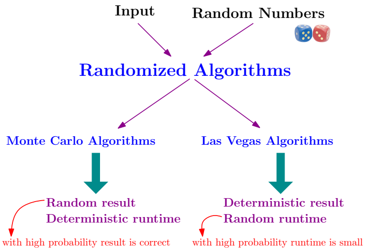
```
k ← 1
while A[k] ≠ 1 do
    k ← RANDOM(1 ··· n)
return k
```

Quality: correct

expected runtime: 4



Las Vegas and Monte Carlo Algorithms



- **Monte Carlo Algorithms**
  - Guaranteed to run in a fixed time
  - Outputs a correct answer with some (high) probability
- **Las Vegas Algorithms**
  - Guaranteed to output the correct answer
  - Running time is random (with high probability runtime is small)

Can always convert a Las-Vegas algorithm into a Monte Carlo algorithm

▷ Stop the algorithm after a certain point

But no method is known for the other way - needs efficient verification