Polynomial Time Reduction

- Polynomial Time Reduction Definition
- Reduction by Equivalence
- Reduction from Special Cases to General Case
- Reduction by Encoding with Gadgets
- Transitivity of Reductions
- Decision, Search and Optimization Problem
- Self-Reducibility

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Problem A is polynomial time reducible to Problem B, $A \leq_{p} B$

If any instance of problem A can be solved using a polynomial amount of computation plus a polynomial number of calls to a solution of problem B

If any algorithm for problem B can be used [called (once or more) with *'clever'* legal inputs] to solve any instance of problem A



Algorithm for A transforms an instance x of A to an instance y of B. Then transforms B(y) to A(x)

Reduction from special case to general case

VERTEX-COVER
$$(G, k) \leq_{p}$$
 SET-COVER (U, S, k')

- Let \mathcal{A} be an algorithm solving SET-COVER (U, \mathcal{S}, k')
- Let [G, k] be an instance of the VERTEX-COVER problem
- Make U = E, k' = k
- $S = \{S_1, \ldots, S_n\}$, where $S_i = \{e \in E \mid e \text{ is incident on } v_i\}$



Theorem: [U, S] has a set cover of size k iff G has a vertex cover of size k

If $\mathcal{A}(U, \mathcal{S}, k') =$ Yes, then output Yes, else output No

Reduction from special case to general case

We get the following reduction very similarly

INDEPENDENT-SET $(G, k) \leq_p$ SET-PACKING(U, S, k')

The following reductions are even more straight forward. They follow from respective definitions of the problems

 $3-\operatorname{SAT}(f) \leq_p \operatorname{SAT}(f')$

SUBSET-SUM $(U, w, C) \leq_{p} \text{KNAPSACK}(U, w, v, C')$

Please complete their details. Explicitly and formally writing them will help understand the important notion of reduction