

STREAMING ALGORITHMS

- Streaming Model of Computation
- Streaming Algorithms and DFA
- Stream: Motivation and Applications
- Synopsis: Sliding Window, Histogram, Wavelets
- Sampling from Stream: Reservoir Sampling
- Linear Sketch
- Count-Min Sketch
- AMS Sketch

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DFA and Streaming Algorithms

Streaming Algorithms can simulate any DFA using logarithmic bits

If L is a language recognized by a DFA D on at most m states, then L is recognized by a streaming algorithm using at most $\log m$ bits of space

We give a streaming algorithm that simulates $D = (Q, \Sigma, q_0, \delta, F)$

Algorithm Streaming Algorithm to simulate $D =$ on string $w = w_1 w_2 \dots w_n$

$q \leftarrow q_0$

▷ $|Q| = m \implies q$ is a $\log m$ -bit integer

for $i = 1 \rightarrow n$ **do**

$q \leftarrow \delta(q, w_i)$

▷ δ function can be provided as a lookup table

if $q \in F$ **then**

▷ Search for q in F

Accept

else

Reject

Runtime of the algorithm depends on the data structure of δ and F but space consumed is one integer ($\log m$ bits)

Can DFA simulate Streaming Algorithms?

Clearly, No! \exists non-regular languages recognizable by streaming algorithms

However, DFA's are not totally powerless!

For a language $L \subseteq \Sigma^*$ and an integer $n \geq 0$, let

- $L_n = \{w \in L : |w| = n\}$
- $L_{\leq n} = \{w \in L : |w| \leq n\}$

$$L = \{0, 10, 100, 110, 0110, 0100, 1010, 1000, 1110\}$$

- | | | |
|------------------------|----------------------------|--------------------------------------|
| ■ $L_1 = \{0\}$ | ■ $L_2 = \{10\}$ | ■ $L_3 = \{100, 110\}$ |
| ■ $L_{\leq 1} = \{0\}$ | ■ $L_{\leq 2} = \{0, 10\}$ | ■ $L_{\leq 3} = \{0, 10, 100, 110\}$ |

Can DFA simulate Streaming Algorithms?

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However, DFA's are not totally powerless!

Theorem: If $L \subseteq \Sigma^*$ is recognized by a streaming algorithm A using $f(n)$ bits of space on any string of length at most n . Then for every n there is a DFA D with at most $2^{f(n)}$ states, with $L(D)_{\leq n} = L_{\leq n}$

We construct a DFA $D = (Q, \Sigma, q_0, \delta, F)$ from A

- Q : all possible $2^{f(n)}$ values of A 's memory
- Σ : alphabet of A or L
- q_0 : initial value of A 's memory
- δ : mimics how A 's memory changes on σ
- F : states corresponding to A 's memory values where A accepts at the end of string

DFA simulating Streaming algorithm to recognize $1\text{-DOMINANT}_{<2}$

Algorithm Streaming Algorithm A
for $1\text{-DOMINANT}(w = w_1 w_2 \dots w_n)$

$C \leftarrow 0$

$x \leftarrow 0$ $\triangleright C$ is int, x is a bit

for $i = 1 \rightarrow n$ **do**

if $C = 0$ **then**

$C \leftarrow 1$

$x \leftarrow w_i$

else if $C \neq 0$ **AND** $x = w_i$ **then**

$C \leftarrow C + 1$

else

$C \leftarrow C - 1$

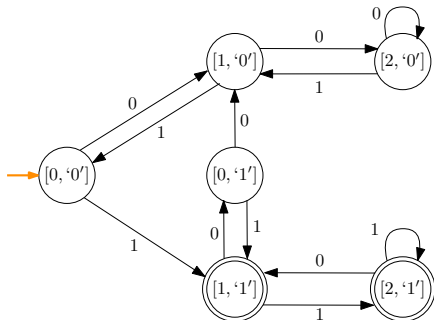
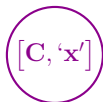
if $C > 0$ **AND** $x = 1$ **then**

Accept

else

Reject

Memory values of A
States of DFA



This DFA agrees with A on strings in L of length ≤ 2