## Algorithms

## Intractable Problems

- Clique
- Independent Set

■ Vertex Cover
■ Set Cover

- Set Packing
- Satisfiability Problem

■ Hamiltonian Cycle and Path

- Traveling Salesman Problem

■ Graph Coloring

- Circuit Satisfiability

■ Knapsack

- Subset Sum
- Prime and Factor
- Partition


## knapscak and subset-sum Problem

■ Given a set $U=\left\{a_{1}, a_{2}, \ldots, a_{n}\right\}$ of objects

- A weight function $w: U \rightarrow \mathbb{Z}^{+}$ $\triangleright w_{i}=w\left(a_{i}\right)$
- A value function $v: U \rightarrow \mathbb{R}^{+}$
$\triangleright v_{i}=v\left(a_{i}\right)$
- And a positive integer $C$

Knapsack $(U, w, v, C, k)$ problem:
Is there a $S \subset U$ such that $\sum_{a_{i} \in S} w_{i} \leq C$ and $\sum_{a_{i} \in S} v_{i}=k$ ?
$\operatorname{SUBSET-SUM}(U, w, C)$ problem: Is there a $S \subset U$ such that $\sum_{a_{i} \in S} w_{i}=C ?$

## Number Theory Problems

$\operatorname{PRIME}(n)$ problem: Is the integer $n$ a prime?

COMPOSITE ( $n$ ) problem: Is the integer $n$ a composite number?
$\operatorname{FACTOR}(n, k)$ problem: Is there a factor $d$ of $n$ such that $2 \leq d \leq k$ ?

These problems are the building blocks of public key cryptography

## The Partition Problem

■ Given a set $U=\left\{a_{1}, a_{2}, \ldots, a_{n}\right\}$ of $n$ positive integers

- Partition $U$ into two subsets $U_{1}$ and $U_{2}$
- $\sum_{a \in U_{1}} a=\sum_{a \in U_{2}} a$
$\triangleright$ balanced partition
- Also called the number bipartition problem

■ If $\left|\sum_{a \in U_{1}} a-\sum_{a \in U_{2}} a\right|=k$, then it is called $k$-imbalanced bipartition PARTITION $(U, k)$ problem: Is there a $k$-imbalanced biparition of $U$ ?

Childhood team selection actually was a greedy algorithm for bipartition
Multiway partition problem is an interesting extension of this problem

## Partition Problem: Applications

Multiprocessor Scheduling
Assign tasks to two identical processors to minimize the MAKESPAN
MAKESPAN is the latest finishing processor

- $U$ is the set of tasks
- A balanced bipartition of $U$ (tasks in each part to be run on one processor) minimizes the makespan

For $k$ processors, it becomes the $k$-way partition problem

## Partition Problem: Applications

## Scoring Based Voting

Three candidates $A, B, C$ and voters with weighted votes. Each voter votes to veto a candidate. The candidate with the smallest total weight of vetoes wins.

We have a subset of $n$ voters with weights $a_{1}, a_{2}, \ldots, a_{n}$, who wants to select candidate $A$. How should they cast their vetoes to ensure $A$ wins.

■ $U=\left\{a_{1}, a_{2}, \ldots, a_{n}\right\}$
■ U wants to elect $A$, they should "bipartition" their vetoes for $B$ and $C$
■ This will maximize the minimum vetoes for $B$ and $C$

