# Dynamic Programming

- (Weighted) Independent Set in Graphs
- Weighted Independent Sets in Path
- Dynamic Programming Formulation
- Implementation and Backtracking

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# The Path Graph

The path graph is a connected graph with two nodes of degree 1 and the other n - 2 vertices of degree 2



▷ So a path is a tree

**Input:** A node weighted graph  $G = (V, E), w w : V \to \mathbb{R}^+$ 

**Output:** An independent set of G of maximum cardinality weight

A company wants to open restaurants on the motorway

- Designated service areas  $s_1, \ldots, s_n$  every 7 kilometers
- A restaurant at  $s_i$  gives estimated profit  $p_i$
- No two restaurants can be located within 10 km of each other

#### Select a subset of sites to maximize total profit

Problem can be modeled by a node weighted path graph

- Each site  $s_i$  is a vertex with weight equal to  $p_i$
- If two sites are within 10 km of each other make an edge between the corresponding vertices
  ▷ note: we get a path graph

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No consecutive vertices can be chosen



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### Greedy Approach:

- Select a node with max weight
- Mark its neighbors as incompatible
- Repeat the process with remaining unmarked nodes



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### Divide & Conquer approach-1:

- Divide *P* into left and right halves
- Find max weight independent sets in both
- Combine the two sets to get the answer



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