## Algorithms

## Huffman Code

- Data Compression
- Lossy and Lossless Compression
- Adaptive and non-Adaptive Compression
- Fixed and Variable length Codes

■ Prefix Free Codes

- Binary Tree Representation
- Goodness Measure
- Generic Greedy Algorithm

■ Huffman Code
■ Optimality and Implementation

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## Data Compression

Data Compression is used in many computer science areas for reduced
■ Computational complexity of data processing

- Storage complexity
- Communication complexity


## Compression Scheme

Let $\mathcal{D}$ be the set of all possible documents (files/input) and $\mathcal{D}^{\prime}$ is set of all possible output documents

A compression scheme has two algorithms

- A compression algorithm $f: \mathcal{D} \rightarrow \mathcal{D}^{\prime}, \quad f(x)=y$

■ A decompression algorithm $g: \mathcal{D}^{\prime} \rightarrow \mathcal{D}, \quad g(y)=x^{\prime}$

## Lossy and Lossless compression

A compression scheme is

- Lossless if $g(y)=f^{-1}(y)=x$ for all $x \in \mathcal{D}$ such that $f(x)=y$

■ Used in Huffman code, .gif, .png

■ Lossy if $g(y) \sim x$ for all $x \in \mathcal{D}$ such that $f(x)=y$

- Similarity between $g(y)$ and $x$ is measured by some error function

■ Used in .mp3, .mpg, .jpg

## Data Compression

source: percona.com data

Lossless Compression Huffman
Shannon Fano
Lossy Compression JPEG MPEG


$\simeq$ original data


## Adaptive and non-adaptive

A Compression scheme can be

- Non-adaptive: - Assumes prior knowledge of the data
- e.g. ' $e$ ' is the most common character in English language documents
- 'the' is the most common word
- Adaptive: - Assumes no prior knowledge of the data
- can build such knowledge (e.g. frequencies in the input document)
- this knowledge will be adaptive to the actual document


## Binary codes

A binary code is a compression scheme with $\mathcal{D}^{\prime}$ as bit-strings

Suppose a file $D$ is 100,000 characters long
$D$ has only 6 unique characters (symbols)
Frequencies of each character is as follows

Characters and their frequencies in $F$

|  | a | b | c | d | e | f | Total |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | ---: |
| Frequency in 000's | 45 | 13 | 12 | 16 | 9 | 5 | 100 |

Find a binary code that encodes $D$ using minimum number of bits

## Fixed versus Variable length codes

## Fixed Length Code

- Fixed number of bits for each symbol (character)

■ e.g. ASCII (7 bits) and Unicode (UTF-8, UTF-16)

- ASCII can represent $2^{7}=128$ symbols


## Variable Length Code

- Variable number of bits for each symbol

■ Can use fewer bits for more frequent symbols

- e.g. Huffman code

■ Difficult to find, needs compression scheme

## Fixed versus Variable length codes

| Characters | a | b | c | d | e | f |  | Total |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Frequency | 45k | 13k | 12k | 16k | 9 k | 5 k |  | 100k |
| Fixed-Length Code | 000 | 001 | 010 | 011 | 100 | 101 |  | 100k |
| Variable Length Code | 0 | 101 | 100 | 111 | 1101 | 1100 |  | 224 k |

Variable length code uses about $25 \%$ less space

## Fixed versus Variable length codes

| Characters | a | b | c | d |
| :--- | :---: | :---: | :---: | :---: |
| Frequency | $\|\mid$ | 5 | 3 | 1 |
|  |  |  |  |  |
| Fixed-Length Code | $\|\mid$ | 00 | 01 | 10 |
| Variable Length Code | 0 | 0 | 10 | 110 |

Let the string be a abbaabl $\mathbf{c} \mathbf{d}$

- Encoding under the fixed length code and it's length is $00000101000000011011 \rightarrow 2 \times 10=20$ bits

■ Encoding under the variable length code and it's length is $00101000010110111 \rightarrow 1 \cdot 5+2 \cdot 3+3 \cdot 1+3 \cdot 1=17$ bits

