Basic Graph Algorithms

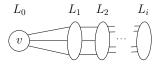
- Exploring Graphs
- Depth First Search
- DFS Forest Start and Finish Time
- DAG, Topological Sorting
- Strongly Connected Components
- Breadth First Search
- Bipartite Graphs

Imdad ullah Khan

Breadth First Search (BFS)

- Breadth First Search is an alternative algorithm to explore graphs
- Starting from a vertex v, BFS explores all vertices that are reachable from v
- BFS first explores neighbors of v, then neighbors of neighbors of v and so on

Starting from ν , explore outwards, adding nodes to one "layer" at a time



- $L_0 := \{v\}$
- $L_1 :=$ neighbors of L_0
- $L_2 :=$ those unvisited nodes having an edge from L_1
- *L*_{*i*+1} := those unvisited nodes having an edge from *L*_{*i*}

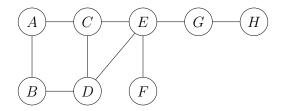
Algorithm BFS(G)

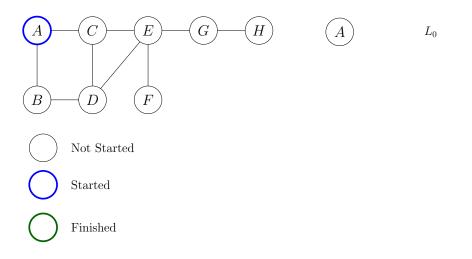
 $visited \leftarrow ZEROS(n)$
for $v \in V$ do
if visited[v] = 0 then
BFS-EXPLORE(v)

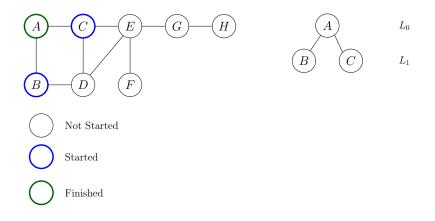
Algorithm BFS-EXPLORE(*s*)

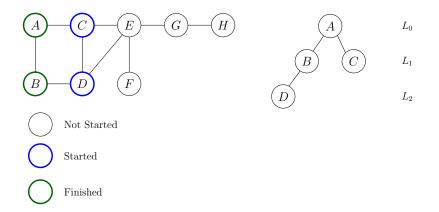
 $\begin{array}{l} \textit{visited}[v] \leftarrow 1 \\ \texttt{ENQUEUE}(Q, s) \\ \texttt{while } Q \neq \emptyset \ \texttt{do} \\ v \leftarrow \texttt{DEQUEUE}(Q) \\ \texttt{for } u \in N(v) \ \texttt{do} \\ \texttt{if } \textit{visited}[u] = 0 \ \texttt{then} \\ \textit{visited}[u] \leftarrow 1 \\ \texttt{ENQUEUE}(Q, u) \end{array}$

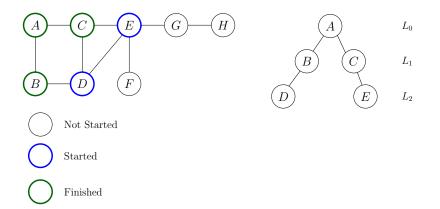
Running time: O(|E| + |V|) = O(n + m)

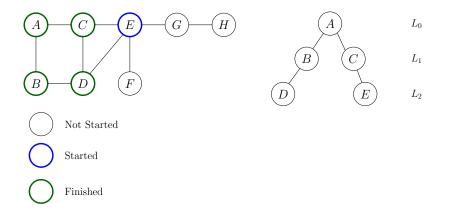


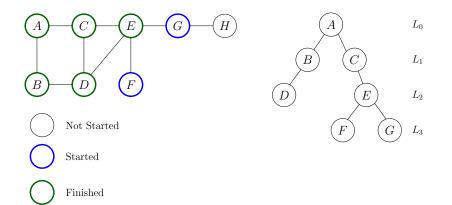


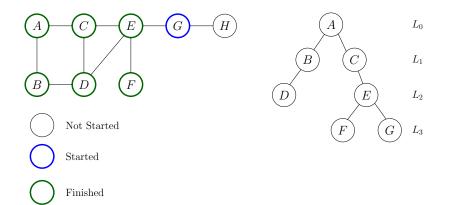


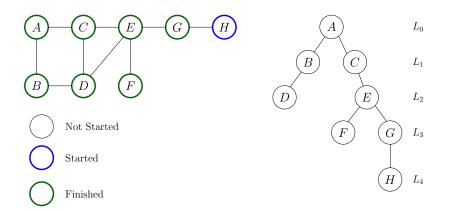


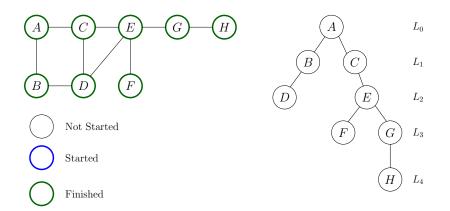


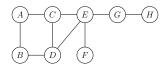




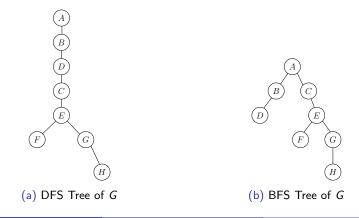








(a) Graph G



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Connected Components via BFS

Find number of connected components in graph G

Algorithm BFS(G)

 $\begin{array}{l} \textit{visited} \leftarrow \textit{ZEROS}(n) \\ \textit{complD} \leftarrow 0 \\ \textit{for } v \in V \textit{ do} \\ \textit{if } \textit{visited}[v] = 0 \textit{ then} \\ \textit{complD} \leftarrow \textit{complD} + 1 \\ \textit{BFS-EXPLORE}(v) \end{array}$

Algorithm BFS-EXPLORE(*s*)

```
\begin{array}{l} \textit{visited}[v] \leftarrow 1 \\ \texttt{ENQUEUE}(Q, s) \\ \texttt{while } Q \neq \emptyset \ \texttt{do} \\ v \leftarrow \texttt{DEQUEUE}(Q) \\ \texttt{for } u \in N(v) \ \texttt{do} \\ \texttt{if } \textit{visited}[u] = 0 \ \texttt{then} \\ \textit{visited}[u] \leftarrow 1 \\ \texttt{ENQUEUE}(Q, u) \end{array}
```

This finds just the number of connected components

 Can be extended to find nodes in all components, components sizes, the largest/smallest components sizes etc.

Running Time: O(|V| + |E|)

Applications of BFS

- Superimpose all graph edges onto the BFS tree
- What can we say about edges?
- (forward edges, tree edges, backward edges, cross edges)
- This observation is the key to many applications of BFS
- Shortest Paths, Bipartite testing, Betweenness Centrality