Basic Graph Algorithms

- Exploring Graphs
- Depth First Search
- DFS Forest Start and Finish Time
- DAG, Topological Sorting
- Strongly Connected Components
- Breadth First Search
- Bipartite Graphs

Imdad ullah Khan

Reachability

Given a graph G = (V, E)

A vertex v is reachable form u, if there exists a path from u to v

• R(u): the set of vertices reachable from u

 $R(u) : \{v : \exists a \text{ path from } u \text{ to } v\}$

Fundamental graph exploration problems

Given s and u, is u ∈ R(s)?
Given s, find R(s)

A more constructive/algorithmic definition of reachability:

v is reachable form s, if v is a neighbor of s or v is reachable from a neighbor of s

$$R(s) = \{s\} \bigcup_{x \in N(s)} R(x)$$

Explore

Input: A graph G = (V, E) and a node $s \in V$ **Output:** All nodes that are reachable from *s*, R(s)

- R(s) is saved as a bit-vector visited[1...n]
- $visited[i] = 1 \Leftrightarrow v_i \in R(s)$

• Populate visited[·] using $R(s) = \{s\} \bigcup_{x \in N(s)} R(x)$

AlgorithmREC-EXPLORE(s)visited[\cdot] \leftarrow ZEROS(n)functionEXPLORE(G,s)visited[s] \leftarrow 1for $x \in N(s)$ doif visited[x] = 0 thenEXPLORE(G,x)

Algorithm ITR-EXPLORE(*s*)

 $visited \leftarrow ZEROS(n)$ INSERT(todo, s)while $todo \neq \emptyset$ do $u \leftarrow REMOVE(todo)$ $visited[u] \leftarrow 1$ for $x \in N(u)$ do $if \ visited[x] = 0 \ then$ INSERT(todo, x)

▷ fixed order

Exploring the whole graph

- EXPLORE is still under-specified to analyze its runtime
- The EXPLORE procedure visits only the portion of graph that is reachable from the given source s
- To examine the rest of the graph, we restart the procedure elsewhere, at some unvisited vertex
- What if we start EXPLORE from some already visited vertex?
- Notice that algorithm readily extends to directed graphs

Input: A graph G = (V, E)**Output:** *'Explore'* the graph *G*

Algorithm DFS(G)

 $visited \leftarrow ZEROS(n)$ for all $s \in V$ do if visited[s] = 0 then DFS-EXPLORE(s) **Algorithm** DFS-EXPLORE(*s*)

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Depth First Search: Forest

Algorithm DFS(G)

 $visited \leftarrow ZEROS(n)$ for all $s \in V$ do if visited[s] = 0 then DFS-EXPLORE(s) **Algorithm** DFS-EXPLORE(*s*)

 $\begin{array}{l} \textbf{function DFS-EXPLORE}(s) \\ \textit{visited}[s] \leftarrow 1 \\ \textbf{for } x \in N(s) \textbf{ do} \\ \textbf{if } \textit{visited}[x] = 0 \textbf{ then} \\ \text{DFS-EXPLORE}(x) \end{array}$

When call to DFS-EXPLORE(s) is executed, all vertices in R(s) are visited

- When DFS-EXPLORE(u) is finished one 'DFS tree' is formed containing all vertices reachable from u
- The next DFS-EXPLORE(v) called from outer loop forms a new tree



Algorithm DFS(G)

 $visited \leftarrow ZEROS(n)$
for all $s \in V$ do
if visited[s] = 0 then
DFS-EXPLORE(s)

Algorithm DFS-EXPLORE(*s*)

function DFS-EXPLORE(s) $visited[s] \leftarrow 1$ for $x \in N(s)$ do if visited[x] = 0 then DFS-EXPLORE(x)

•
$$G = (V, E), |V| = n, |E| = m$$

• Initialization and visited checking takes O(n) time

Total runtime within all EXPLORE calls is

 $\sum_{i=1}^n deg(v_i) = 2m$

• Total runtime is O(n+m)

DFS: Summary

Algorithm DFS(G)

visited $\leftarrow \text{ZEROS}(n)$ for all $s \in V$ do if visited[s] = 0 then DFS-EXPLORE(s)

Algorithm DFS-EXPLORE(*s*)

- DFS explores the entire graph
- Explores one neighbor first (in depth), before going to next neighbor
- Works both for undirected and directed graphs
- Runtime of O(n + m) is equal to reading the input graph
- DFS doesn't add any cost (asymptotically) to any graph algorithm, we typically do it as a pre-processing step
- Answers many questions
 - is graph connected, how many components in the graph, find $R(s), \cdots$
- A fundamental algorithm with many applications; we will discuss some