Induction

- Principle of Mathematical Induction
- Proofs by Induction
- Strong Induction
- Well Ordering Principle

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Principle of Mathematical Induction

 $\left[P(0) \land \forall k \ge 0 \left[P(k) \to P(k+1) \right] \right] \longrightarrow \forall n \ge 0 P(n)$

Proof using Induction

- *Basis Step:* Prove *P*(0) is true
- IH: Assume P(n)
- Inductive Step: Using P(n), prove P(n+1)

Principle of Strong Mathematical Induction

 $\left[P(0) \land \forall k \ge 0 \left[\forall \ 0 \le i \le k \ P(i) \to P(k+1) \right] \right] \longrightarrow \forall n \ge 0 \ P(n)$

Proof using Strong Induction

- *Basis Step:* Prove *P*(0) is true
- *IH*: Assume P(k) is true for all $1 \le k \le n$
- *IS*: Using $\forall k \leq n P(k)$, prove P(n+1)

Principle of Well Ordering

Any nonempty set of nonnegative integers has a smallest element

What is the smallest element of

 ICP 10-11
 $\{x \in \mathbb{Z} : x < 7\}$

 ICP 10-12
 $\{x \in \mathbb{R} : 0 < x < 1\}$

 ICP 10-13
 $\{3.9 + \frac{1}{n} : n \in \mathbb{Z}^+\}$

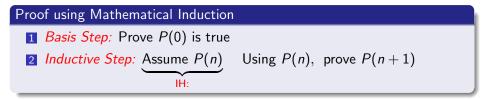
 ICP 10-14
 $\{\lfloor 3.9 + \frac{1}{n} \rfloor : n \in \mathbb{Z}^+\}$

Principle of Well Ordering

Any nonempty set of nonnegative integers has a smallest element

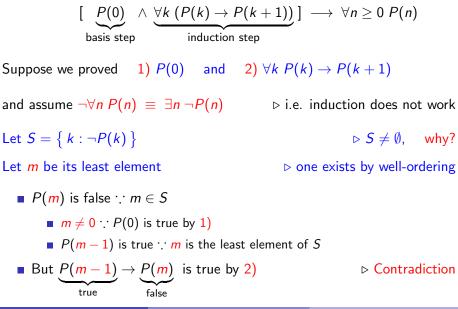
Principle of Mathematical Induction

 $\left[P(0) \land \forall k \ge 0 \left[P(k) \to P(k+1) \right] \right] \longrightarrow \forall n \ge 0 P(n)$



We argue why induction make sense using the principle of well ordering and how they are essentially equivalent

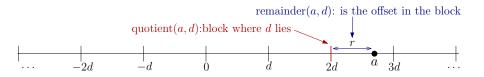
Why induction work



Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

- q is called the quotient
- r is called the remainder
- d is called the divisor
- a is called the dividend



Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

quotient and remainder when 11 is divided by 5

 $11~=~5\times2+1$

Notation:

q = 2 = 11 div 5r = 1 = 11 mod 5

Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

quotient and remainder when -11 is divided by 5

 $-11~=~5\times-3+4$

Notation:

q = -3 = -11 div 5r = 4 = -11 mod 5

Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

$11~=~5\times2+1$

Given positive *a*, repeatedly subtract *d* until what remains *r* is $0 \le r < d$

$-11 = 5 \times -3 + 4$

Given negative *a*, repeatedly add *d* until what remains *r* is $0 \le r < d$

Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

Existence:
$$S = \{a - dq \ge 0 : q \in \mathbb{Z}\}$$
 $S \ne \emptyset$
Let $r = a - dq_0$ be the least element of S $r \ge 0$
 $r < d$?
If $r \ge d$, then $a - d(q_0 + 1) = r - d \ge 0$
 $r - d < r$ is in S , \triangleright a contradiction to minimality of r

There are q and $0 \le r < d$ with a = dq + r

Theorem (The Division Algorithm)

Let **a** be an integer and **d** a positive integer. Then there are unique integers **q** and **r**, with $0 \le r < d$ such that a = dq + r

Uniqueness:

Let $a = dq_1 + r_1$ and $a = dq_2 + r_2$ with $q_1 \neq q_2$ and $0 \leq r_1 \neq r_2 < d$

$$\underbrace{dq_1+r_1}_{a} - \underbrace{dq_2+r_2}_{a} = 0 \implies d(q_1-q_2) = r_2-r_1$$

This means
$$d$$
 divides $(r_2 - r_1) \implies$ either $\begin{cases} r_2 = r_1 & \text{or} \\ d \le |(r_2 - r_1)| \end{cases}$

$$\begin{bmatrix} 0 \le r_1 \neq r_2 < d \end{bmatrix} \implies -d < r_2 - r_1 < d \implies |r_2 - r_1| < d$$

 \implies $r_1 = r_2$

From $r_1 = r_2$, we also get $q_1 = q_2$