Predicate Logic

- Predicates and Propositional Functions
- Universal and Existential Quantifiers
- Negating Quantified Statements
- Nested Quantified Expressions

Imdad ullah Khan

Quantified Expression: Recap

- Propositional function becomes proposition when specific value is given to variable
- Quantifiers make it proposition for a range of values
- Universal Quantifier: $\forall \forall x P(x) := P(x)$ (is true) for all values of x in the UoD

Proposition $\forall x P(x)$ is **True** iff for every x in UoD, P(x) is **True**

■ Existential Quantifier: ∃

 $\exists x P(x) := P(x)$ (is true) for some value(s) of x in the UoD

Proposition $\exists x P(x)$ is **True** iff for at least one x in UoD, P(x) is **True**

Predicates can have more than one variables

- P(x, y) : x teaches course y in LUMS
- Q(x, y, z): x is an instructor of course y in university z
- S(x, y, z) : x + y = z

P(Pythogarus, CS210) = ?P(Pythogarus, y) = ?P(x, CS210) = ?

Each variable needs to be given value or quantified to make the predicate a proposition

Nested Quantified Expressions: Binding Variables

Predicates can have more than one variables

Each variable needs to be given value or quantified to make the predicate a proposition

• $\forall x P(x, y)$ is not a proposition

▷ The variable x is 'bound', while y is 'free'

• $\forall x \forall y P(x, y)$ is a proposition

▷ both variables are bound

• $\forall x \exists y \forall z \ Q(x, y, z)$ is a proposition

> all variables are bound

Nested Quantified Expressions: Binding Variables

$$\forall x, \forall y \ P(x, y) \equiv^? \quad \forall x \ P(x, y) \land \forall y \ P(x, y)$$

- LHS: Both x and y are bound
- RHS: x is bound in the first predicate and y is free
- RHS: y is bound in the second predicate and x is free
- RHS is not even a proposition

Nested Quantified Expressions: Examples

Translate these statements to logical expressions with quantifiers The sum of two positive integers is always positive

 $\forall x \forall y ((x > 0) \land (y > 0)) \rightarrow (x + y > 0)$

Every real number except zero has a multiplicative inverse

 $\forall x \exists y \ (x \neq 0) \rightarrow (xy = 1)$

If a person is female and is a parent, then this person is someone's mother $\forall x \exists y \ (F(x) \land P(x)) \rightarrow M(x, y)$

Everyone has exactly one best friend

 $\forall x \exists y \forall z \ (B(x,y) \land (y \neq z)) \rightarrow \neg B(x,z)$

Nested Quantified Expressions: Examples

Translate the following logical expressions into English statements

 $\forall x \exists y \ F(x,y) \land x \neq y$

Everyone has at least one friend

 $\exists y \forall x \ F(x,y) \land x \neq y$

There is at least one person who is friend of everyone

Are these statements logically equivalent?

Nested Quantified Expressions: Examples

Translate the logical expressions into English statements

 $\forall x \exists y \ (x \neq 0) \rightarrow (xy = 1)$

Every real number except zero has a multiplicative inverse

$$\exists y \forall x \ (x \neq 0) \rightarrow (xy = 1)$$

Some real number is multiplicative inverse of every non-zero real number

Are these statements logically equivalent?

Order of quantifiers is extremely important

 $\forall x \exists y \ P(x, y)$ is not the same as $\exists y \forall x \ P(x, y)$

P(x, y) : x teaches course y in LUMS

- $\forall x \exists y P(x, y)$: For every instructor, there is a course that (s)he teaches
- $\exists y \forall x P(x, y)$: There is a course that every instructor teaches
- $\exists x \forall y P(x, y)$: There is an instructor who teaches all the courses
- $\forall y \exists x P(x, y)$: For every course, there is an instructor who teaches it

Order of quantifier is extremely important

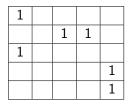
 $\forall x \exists y \ P(x,y)$ is not the same as $\exists y \forall x \ P(x,y)$

Q(x, y, z): x is an instructor of course y in university z

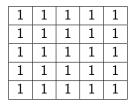
• $\forall z \ \forall y \ \exists x \ Q(x, y, z)$: In every uni. for every course there is an instructor

• $\forall z \exists x \forall y Q(x, y, z)$: In every uni. there is an instructor for all courses

■ $\exists x \forall z \forall y Q(x, y, z)$: There is an instructor for every course in every uni.









- $\forall \text{ row } x \exists \text{ column } y A[x][y] = 1$
- Every row has at least one 1
- Both tables satisfy this

- \exists column $y \forall$ row x A[x][y] = 1
- There is a column with all 1's
- Table 1 does not satisfy this
- Table 2 satisfy this

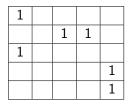


Table: 1

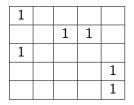
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1
1	1	1	1	1



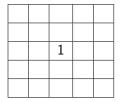
- $\forall \text{ row } x \forall \text{ column } y A[x][y] = 1$
- Every cell has a 1
- Table 1 does not satisfy this
- Table 2 satisfy this

- \forall column $y \forall$ row x A[x][y] = 1
- Every cell has a 1
- Table 1 does not satisfy this
- Table 2 satisfy this

Nested Quantified Expressions: Order of Quantifiers









- $\blacksquare \exists row x \exists column y A[x][y] = 1$
- There is a 1 in the table
- Both tables satisfy this

- \exists column $x \exists$ row y A[x][y] = 1
- There is a 1 in the table
- Both tables satisfy this

Order of quantifier is extremely important

$\forall x \exists y \ P(x,y)$ is not the same as $\exists y \forall x \ P(x,y)$



If both quantifiers are the same then order does not matter

$$\forall x \forall y P(x,y) \equiv \forall y \forall x P(x,y) \text{ and } \exists x \exists y P(x,y) \equiv \exists y \exists x P(x,y)$$

Goldbach conjecture (1742)

Every even integer greater than 2 can be written as sum of two primes

$$orall n \in \mathbb{N}, ((n > 2) \land (n ext{ even}))
ightarrow (\exists p, q \in \mathbb{P}, n = p + q)$$

Vinogradov (1937)

Every sufficiently large odd number is the sum of three primes

$$\exists K \in \mathbb{N}, \forall n \geq K, ((n \text{ odd}) \rightarrow \exists p, q, r \in \mathbb{P}, n = p + q + r)$$

Statement	When True?	When False?	
	P(x, y) is true for every pair x, y	There is a pair x, y for which $P(x, y)$ is false	
$\forall x \exists y \ P(x,y)$	For every x, there is a y for which $P(x, y)$ is true	There is an x such that $P(x, y)$ is false for every y	
$\exists x \forall y \ P(x,y)$	There is an x for which $P(x, y)$ is true for every y	For every x there is a y for which $P(x, y)$ is false	
$\exists x \exists y \ P(x, y) \\ \exists y \exists x \ P(x, y) \end{cases}$	There is a pair x, y for which $P(x, y)$ is true	P(x, y) is false for every pair x, y	

Express each of these system specifications using predicates with UoD's, quantifiers, and logical connectives

ICP 2-27 At least one console must be accessible during every fault condition

A(x, y): Console x is accessible during fault condition y $\forall y \exists x \ A(x, y)$ Express each of these system specifications using predicates with UoD's, quantifiers, and logical connectives

ICP 2-28 The e-mail address of every user can be retrieved whenever the archive contains at least one message sent by every user on the system

R(y): Email of user y can be retrieved M(x, y): Message x has been sent by user y $\forall y \forall z \exists x \ M(x, z) \rightarrow R(y)$ Express each of these system specifications using predicates with UoD's, quantifiers, and logical connectives

ICP 2-29 For every security breach there is at least one mechanism that can detect that breach if and only if there is a process that has not been compromised

D(x, y): Mechanism x can detect breach y

C(z): Process z has been compromised

 $\forall y \exists x \exists z \ D(x, y) \iff \neg C(z)$

Express each of these system specifications using predicates with UoD's, quantifiers, and logical connectives

ICP 2-30 There are at least two paths connecting every two distinct endpoints on the network

C(p, x, y): Path p connects endpoint x and y

 $\forall x \forall y \neq x \exists p_1 \exists p_2 \neq p_1 \ C(p_1, x, y) \land C(p_2, x, y)$

Translate these statements expressing some mathematical fact into English Let the Universe of discourse for all variables be real numbers

ICP 2-32
$$\forall x \exists y \ (x + y = 0)$$
Every real number has an additive inverseICP 2-33 $\exists x \forall y \ (x + y = y)$ There exists an additive identityICP 2-34 $\forall x \exists y \ ((x \neq 0) \rightarrow (xy = 1))$ Every non-zero real number has a multiplicative inverseICP 2-35 $\exists x \forall y \ (xy = y)$

There exists a multiplicative identity

Negating Nested Quantified Expressions

Recall

$$\neg \forall x \ P(x) \equiv \exists x \ \neg P(x)$$

 $\neg \exists x \ P(x) \equiv \forall x \ \neg P(x)$

Negate nested quantified statements using iterative applications of negating (singly) quantified statements

$$\neg \forall x \exists y \ P(x,y) \equiv \exists x \neg \exists y \ P(x,y) \equiv \exists x \forall y \neg P(x,y)$$

$$\neg \exists x \forall y P(x,y) \equiv \forall x \neg \forall y P(x,y) \equiv \forall x \exists y \neg P(x,y)$$

$$\neg \forall x \forall y P(x,y) \equiv \exists x \neg \forall y P(x,y) \equiv \exists x \exists y \neg P(x,y)$$

$$\neg \exists x \exists y P(x,y) \equiv \forall x \neg \exists y P(x,y) \equiv \forall x \forall y \neg P(x,y)$$

Negating Nested Quantified Expressions

P(x, y) : x teaches course y in LUMS

Negate the following statement in plain English

For every course, there is an instructor teaching it

Next translate it into a quantified expression

Negate the quantified expression

Translate the negated quantified expression into plain English and compare with the one you got earlier

Nested Quantified Expressions: Summary

- All variables in multivariable predicates need to be bound
- Each variable can be quantified differently
- Order of different quantifiers is extremely important
- For same quantifiers order does not matter
- It is tricky to translate English statements to nested quantified statement and vice-versa
- Requires extensive practice
- Negating nested quantified statement is just nested applications negating quantified statements